# Practical Evaluation of Protected RNS Scalar Multiplication

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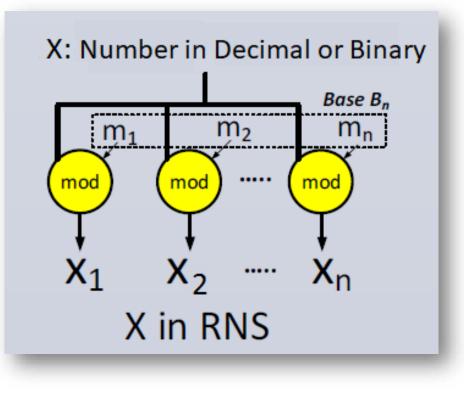




# Outline

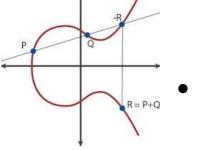
- Residue Number System in Elliptic Curve Cryptography
- Proposed TVLA threshold calculation
- TVLA analysis
- Location and Data Dependent Template Attacks
- Conclusions

#### Residue Number System



X = 50(m1, m2, m3) = (3, 7, 11) (x1, x2, x3) = (2, 1, 6)

# RNS in Elliptic Curve Cryptography



- Elliptic curves defined over prime fields GF(p)
- Modular operations turn easily to RNS modular operations over GF(p)
- RNS mod multiplication usually realized through RNS Montgomery multiplication to avoid modular inversion, but includes base extension
- EC scalar multiplication is the critical operation Q = kP

# LRA Montgomery Power Ladder

Choose base  $B_n$ ,  $B'_n$ . Transform V, R to RNS format using permutation  $p_t$ 

- $R_0 = R$ ,  $R_1 = R + V$ ,  $R_2 = -R$
- Convert  $R_0$ ,  $R_1$ ,  $R_2$  to Montgomery format
- For i= t-1 to 0
  - $R_2 = 2R_2$  in permutation  $p_t$
  - If  $k_i = 1$

 $R_0 = R_0 + R_1$  and  $R_1 = 2R_1$  in permutation  $p_t$ 

else

 $R_1 = R_0 + R_1$  and  $R_0 = 2R_0$  in permutation  $\gamma_t$ 

• Integrity check: if i,k not modified and  $R_0 + V = R_1$  then ret.  $R_0 + R_2$  in permutation  $\gamma_t$ else ret. random value Transform  $R_0 + R_2$  to binary format

# Test Vector Leakage Assessment (TVLA)

- Statistical tests between two trace-sets of acquisition
- Welch's t-test to evaluate if two sets have significant statistical differences

$$s_i = \frac{L_{i,A} - L_{i,B}}{\sqrt{\frac{\sigma_{i,A}^2}{n_A} + \frac{\sigma_{i,B}^2}{n_B}}}$$

• Values above ±4.5, indicates leakage, but TVLA does not exploit it

# t-test Threshold Calibration for TVLA

**Input**  $nt_A$ ,  $nt_B$ : number of traces for groups A,B  $n_s$ : number of samples  $\sigma_A, \sigma_B$ : sampled standard deviation **Output** Threshold value for Welch's t-distribution  $th_t$ 

$$nt_A = nt_B = 4 * 10^3 - 10 * 10^3$$
  
 $n_s = 4 * 10^5 - 8 * 10^5$   
 $\sigma_A = 9.7$ ,  $\sigma_B = 6.1$ 

= +6.3

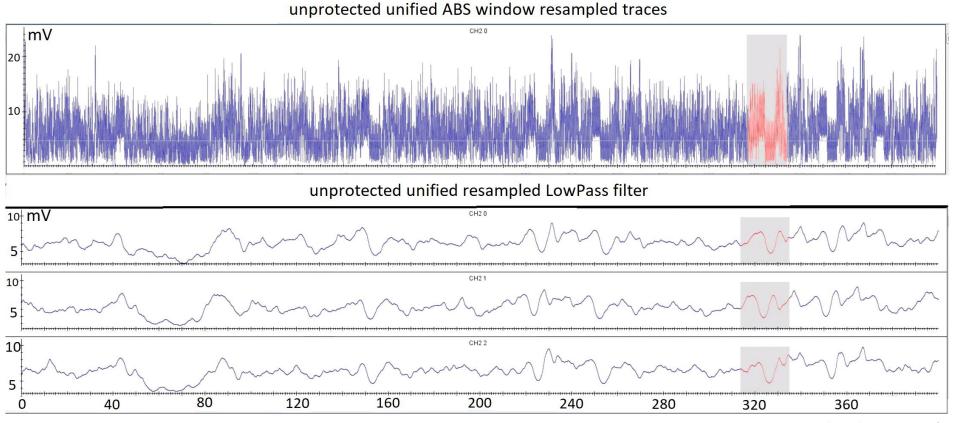
- 1. Choose level of significance  $\alpha$ . Here  $\alpha$ =0.00001
- 2. Family-wise error rate fwer =  $(1 a)^{n_s}$

3. Šidak correction 
$$sidak_{a} = 1 - (1 - a)^{(1/n_s)}$$
  
4. df =  $\left(\frac{\sigma_A^2}{nt_A} + \frac{\sigma_B^2}{nt_B}\right)^2 / \left(\frac{\left(\frac{\sigma_A^2}{nt_A}\right)^2}{nt_A - 1} + \frac{\left(\frac{\sigma_B^2}{nt_B}\right)^2}{nt_B - 1}\right)$   
5. Threshold  $th_t$  =  $|tinv(1 - sidak_a/2, df)|$   $th_t$ 

#### **RNS implementation on BeagleBone**

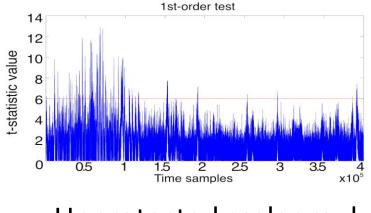
- C Software implementation on ARM Cortex A8
- RNS Montgomery multiplication
- Dedicated and Unified Group Law
- 5 different variations: unprotected, randomized scalar, random input point, random base permutations (LRA), random order of operations

# Processing of Traces – Low Pass Filter

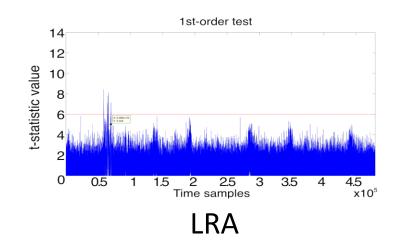


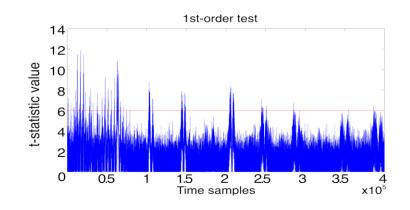
time in µseconds

#### t-test random vs fixed scalar on twisted Edwards curve (a=1, d=2, p= $2^{192} - 2^{64} - 1$ )

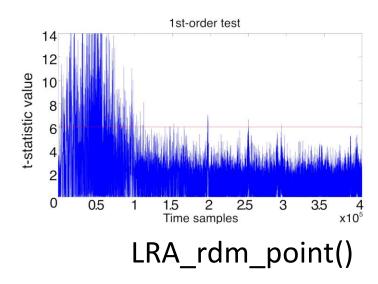


Unprotected scalar mul

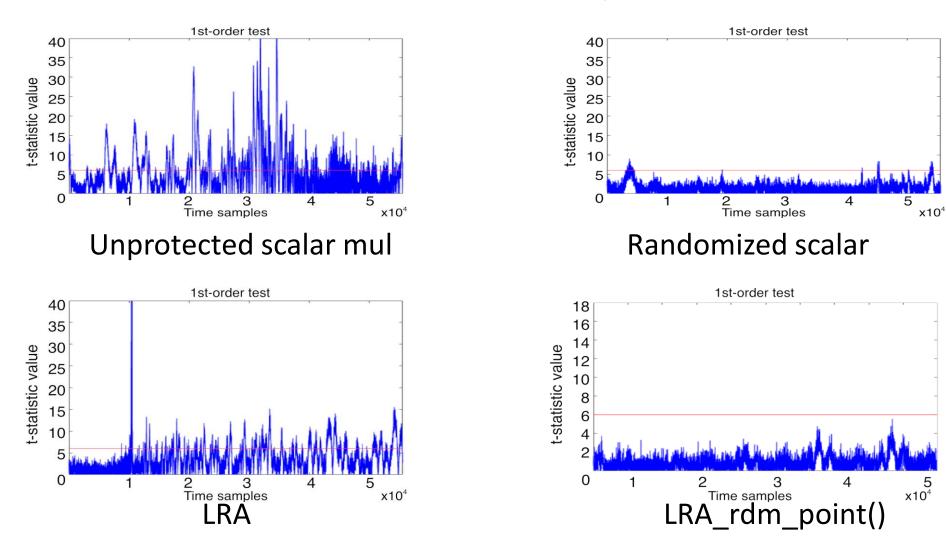




#### Randomized scalar



# t-test random vs fixed point on secure Edwards curve (a=107, d=47, h=4, p= $2^{192} - 2^{64} - 1$ )



# Data Dependent Template Attacks

- The value of a secret variable can be monitored
- Trigger around the key-dependent assignment (if-statement) If k<sub>i</sub> = 1: R<sub>0</sub> = R<sub>0</sub> + R<sub>1</sub> and R<sub>1</sub> = 2R<sub>1</sub> Else: R<sub>1</sub> = R<sub>0</sub> + R<sub>1</sub> and R<sub>0</sub> = 2R<sub>0</sub>
- After alignment, 20k traces. Used half for templates, half for classification
- Success rate 90-91% for the unprotected case, 82-97% for LRA countermeasure activated
- Scalar randomization (65-72%) and LRA randomized RNS operations (55-58%) are good countermeasures

# **Location Dependent Template Attacks**

- Templates created for storage structure that handles the keydependent instruction (doubling) If  $k_i = 1$ :  $R_0 = R_0 + R_1$  and  $R_1 = 2R_1$ Else:  $R_1 = R_0 + R_1$  and  $R_0 = 2R_0$
- Template classification: 95-99.9%
- LRA with randomized operations: 70-83%

# Location Dependent Leakage

- Registers are not really single registers, RNS values are stored in 50-bit chunks result of doubling is stored in different memory locations
- Location dependent leakage was not an expected result
- The normal distributions for  $k_i = 0$  and  $k_i = 1$  for every variation of the implementation are very different (N(-24.3, 9, 7), N(19.6, 6.1))
- Leaky platform capacitors next to each other
- Scalar randomization not an efficient countermeasure
- LRA with randomized operations makes template attacks harder

### **Evaluation Table**

Algorithm	Welch t-test	Welch t-test	TA	TA	PO
	r-vs-f scalar	r-vs-f point	Data	Location	
unprotected	×	×	×	×	0%
rdm_point	×	×	×	×	52%
LRA	×	×	×	×	50%
protected_LRA	×	1	×	×	110%
unprotect_rdm_scal	1	N/A	1	×	19%
rdm_point_rdm_scal	1	N/A	1	×	54%
LRA_rdm_scalar	1	N/A	1	×	51%
protected_rdm_scal	✓	N/A	1	1	110%
unprotect_unified	×	×	X	×	19%
rdm_point_unified	×	×	×	×	99%
LRA_unified	×	×	X	×	72%
protected_unified	1	1	×	×	144%
LRA_nc_rdm_operat	×	1	1	1	76%
LRA_nc_rdm_operat					
_rdm_scalar	✓	N/A	1	1	76%

Pass t-test/secure against templates

Fail t-test/not secure against templates

# Conclusions

- TVLA bounds not rigid; compute according to distribution of traces, number of samples, number of traces
- Randomization of scalar, input point, regularity of MPL are good countermeasures but not enough to avoid leakage
- Different RNS representations do not lower the template success rates
- Randomization of RNS operations protects against templates and less expensive compared to randomization of input point
- Classification using ML algorithms
- Evaluation on an FPGA would give further insights in the security of RNS

#### **THANK YOU FOR YOUR ATTENTION !**



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